

Live Update: The Making of

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If We Functioned Like Computers. . .



*"If you think database patching is onerous,
then try patching a SCADA system
that's running a power plant."*

Kelly Jackson Higgins on the SCADA patch problem, 2013

Solution 1: "Spare time for downtime"



*"In one of the biggest computer errors in banking history,
Chemical Bank mistakenly deducted about \$15 million
from more than 100,000 customers' accounts."*

Saul Hansell, New York Times, 1994

Solution 2: "Roll your upgrades"



"Our research shows that 75% of successful attacks occur against previously known vulnerabilities for which a remediation was already available."

Neil MacDonald, Gartner Research, 2012

Solution 3: "Don't patch, don't tell"



*"All problems in computer science can be solved
by another level of indirection—but that
will usually create another problem."*

Butler Lampson, quoting David Wheeler

Our solution: "Live update"



Live Update in the Real World



Servers protected with Ksplice Uptrack:
100,000+ at more than **700** companies

Updates applied on production systems:
More than **2 million** and counting

How it works



Your Linux vendor releases an update.



Ksplice converts the update into a rebootless update.



You install the update seamlessly, without rebooting.

Source: <http://www.ksplice.com>



Are We There Yet?

```
--- a/drivers/md/dm-crypt.c
+++ b/drivers/md/dm-crypt.c
@@ -690,6 +690,8 @@ bad3:
    bad2:
        crypto_free_tfm(tfm);
    bad1:
        /* Must zero key material before freeing */
+   memset(cc, 0, sizeof(*cc) + cc->key_size * sizeof(u8));
    kfree(cc);
    return -EINVAL;
}
@@ -706,6 +708,9 @@ static void crypt_dtr(...)
    cc->iv_gen_ops->dtr(cc);
    crypto_free_tfm(cc->tfm);
    dm_put_device(ti, cc->dev);
+
+   /* Must zero key material before freeing */
+   memset(cc, 0, sizeof(*cc) + cc->key_size * sizeof(u8));
    kfree(cc);
}
```

Linux kernel security patch for CVE-2006-0095



Existing Live Update Solutions for C

Safe update state

- Update-agnostic characterization, e.g., no updates to active code.
- Problem: extensive patch inspection required for update safety.

State transfer

- Automatic generation of basic type transformers.
- Problem: significant programming effort for complex updates.

Live update mechanisms

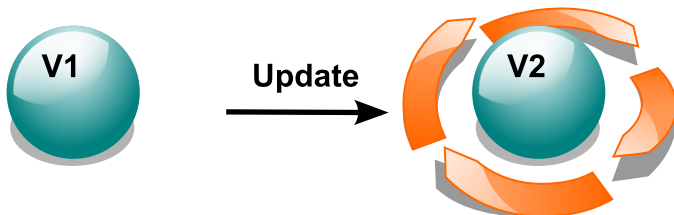
- In-place “*hot patching*” update strategy.
- Problem: *unstable* live update process.



WWW: What We Want

- Support for simple and complex updates of different natures.
- Safe and predictable live update process.
- Automated state transfer and state checking.
- Automatic error recovery (*hot rollback*).
- Stable live update process.

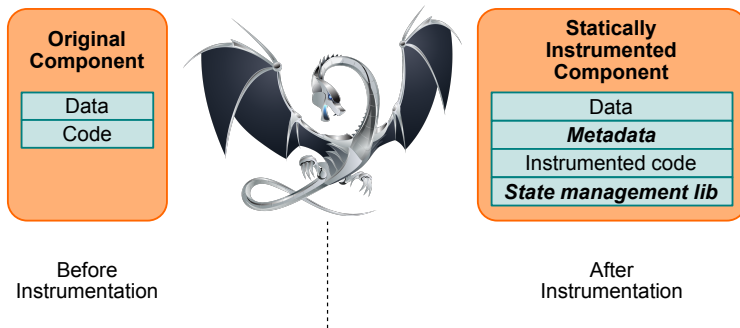




Process-level updates

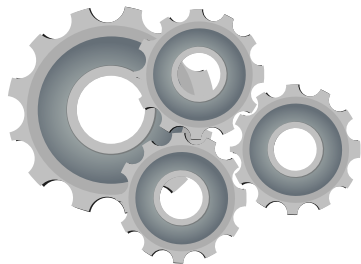


Our Live Update Design



Compiler-based state instrumentation

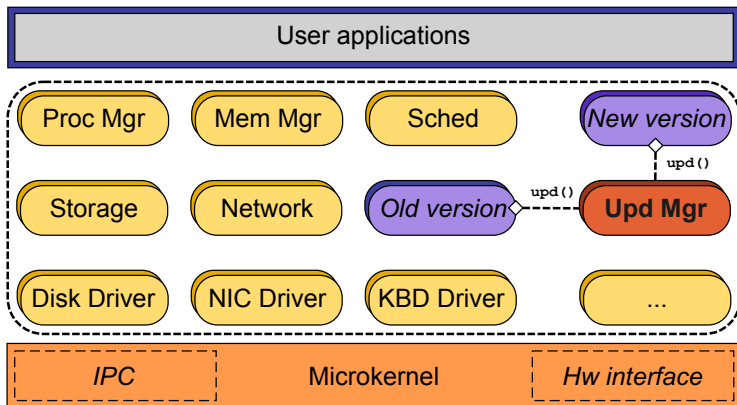




Controlled live update transaction



PROTEOS Architecture



```
static int my_init() {  
    ... //initialization code  
    return 0;  
}  
  
int main() {  
    event_eh_t my_ehs = {init : my_init};  
    sys_startup(&my_ehs);  
    while(1) { // event loop  
        msg_t m;  
        sys_receive(&m);  
        process_msg(&m);  
    }  
    return 0;  
}
```

Event-driven model



```
static int my_init() {  
    ... //initialization code  
    return 0;  
}  
  
int main() {  
    event_eh_t my_ehs = {init : my_init};  
    sys_startup(&my_ehs);  
    while(1) { // event loop  
        msg_t m;  
        sys_receive(&m);  
        process_msg(&m);  
    }  
    return 0;  
}
```

Entry point




```
static int my_init() {  
    ... //initialization code  
    return 0;  
}  
  
int main() {  
    event_eh_t my_ehs = {init : my_init};  
    sys_startup(&my_ehs);  
    while(1) { // event loop  
        msg_t m;  
        sys_receive(&m);  
        process_msg(&m);  
    }  
    return 0;  
}
```

Update point



Live Update Example

```
% prctl mupdate net /bin/net.new \  
-state 'num_pending_writes == 0'
```

```
% prctl mupdate e1000 /bin/e1000.new
```

```
% prctl mupdate-start
```

UM: Live update requested for *net*, *e1000*.

UM: Loading /bin/net.new in memory...

UM: Loading /bin/e1000.new in memory...

UM: Applying changes...

UM: Cleaning up old version...

UM: Live update done.

Multi-component live update



Live Update Example

```
% prctl mupdate net /bin/net.new \  
    -state 'num pending writes == 0'  
  
% prctl mupdate e1000 /bin/e1000.new  
  
% prctl mupdate-start  
  
UM: Live update requested for net, e1000.  
UM: Loading /bin/net.new in memory...  
UM: Loading /bin/e1000.new in memory...  
UM: Applying changes...  
UM: Cleaning up old version...  
UM: Live update done.
```

State filter



Live Update Example

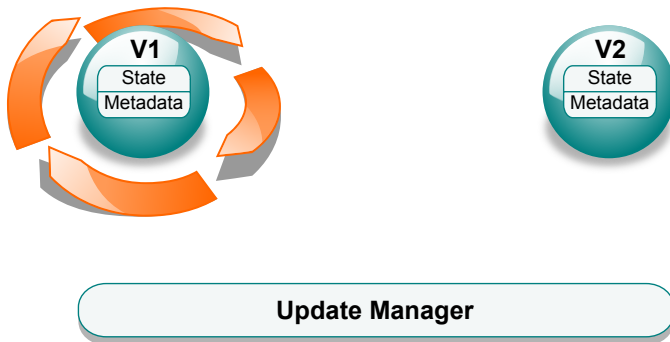
```
% prctl mupdate net /bin/net.new \  
    -state 'num_pending_writes == 0'  
  
% prctl mupdate e1000 /bin/e1000.new  
  
% prctl mupdate-start
```

```
UM: Live update requested for net, e1000.  
UM: Loading /bin/net.new in memory...  
UM: Loading /bin/e1000.new in memory...  
UM: Applying changes...  
UM: Cleaning up old version...  
UM: Live update done.
```

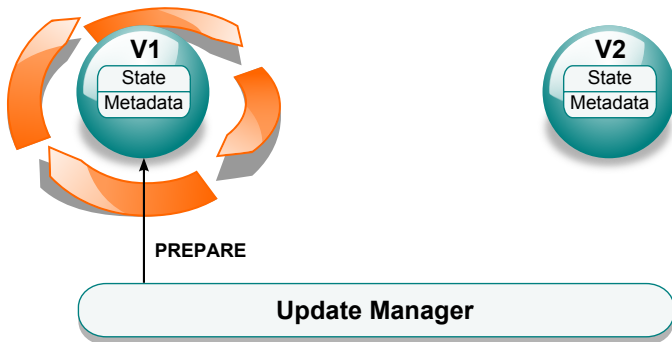
Changes applied automatically



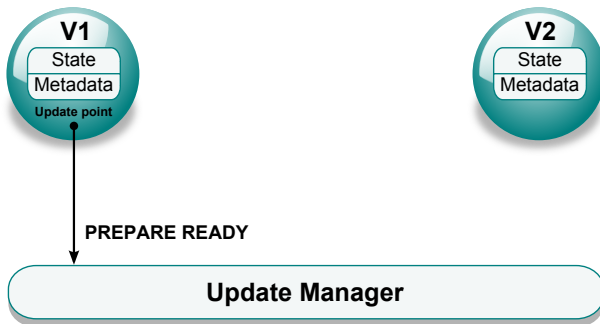
The Live Update Process



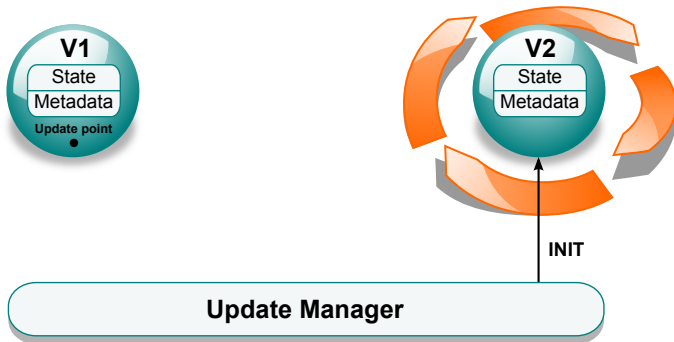
The Live Update Process



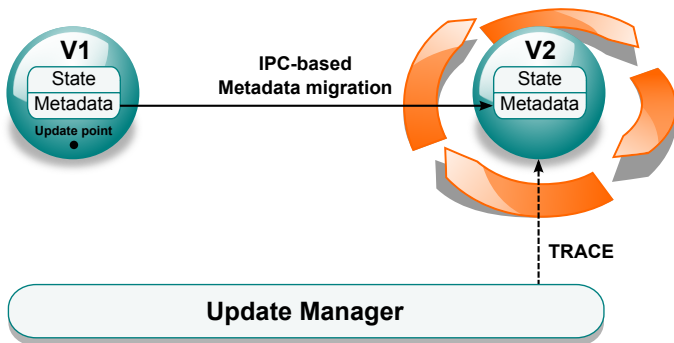
The Live Update Process



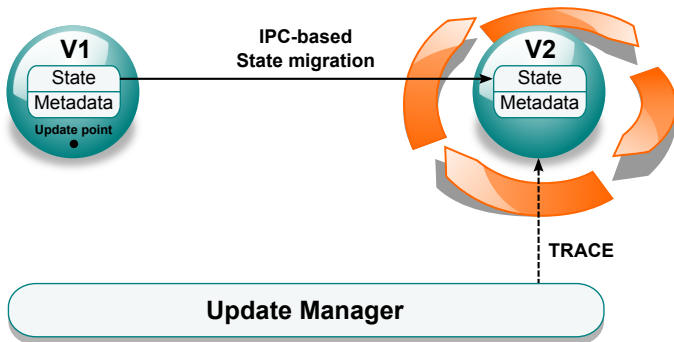
The Live Update Process



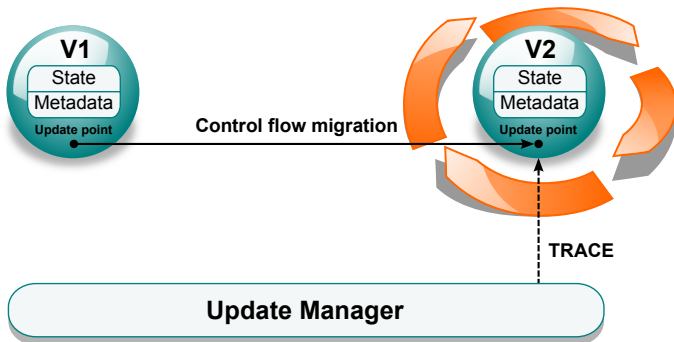
The Live Update Process



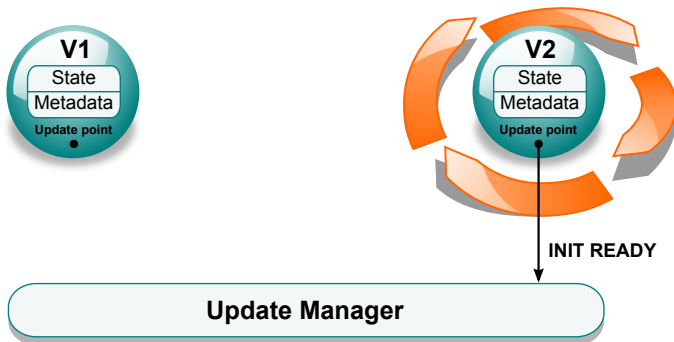
The Live Update Process



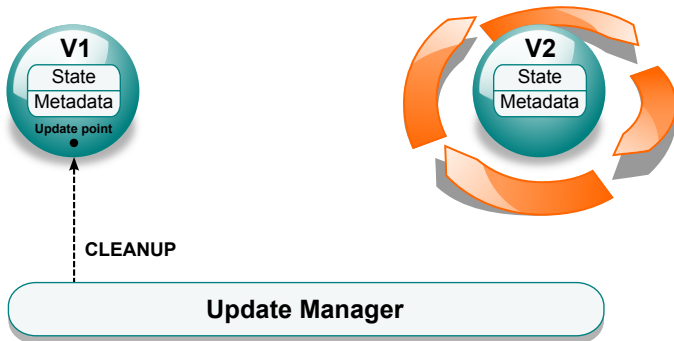
The Live Update Process



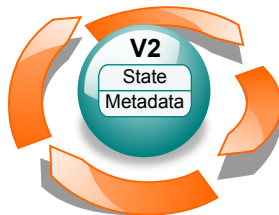
The Live Update Process



The Live Update Process



The Live Update Process



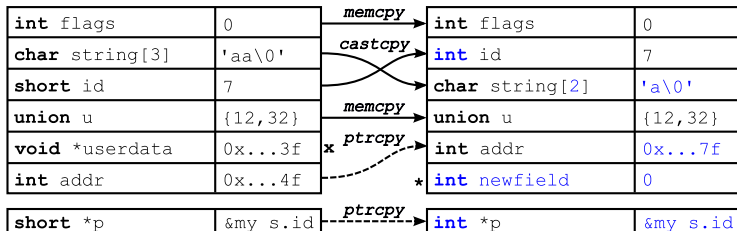
Update Manager



Transfer Strategy

```
struct s { //old version
    int flags;
    char string[3];
    short id;
    union IXFER(my_u) u;
    void *userdata;
    PXFER(int) addr;
} my_s;
short *p = &my_s.id;
```

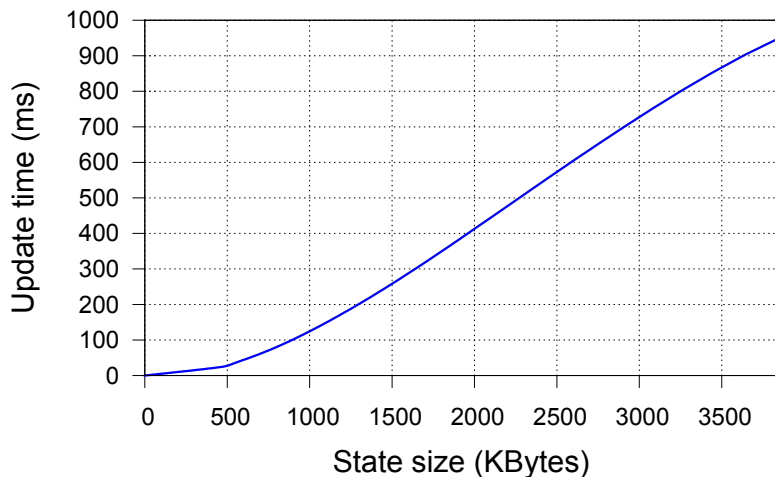
```
struct s { //new version
    int flags;
    int id;
    char string[2];
    union IXFER(my_u) u;
    PXFER(int) addr;
    int newfield;
} my_s;
int *p = &my_s.id;
```



- Applied 50 real updates (~ 15000 LOC) with only 265 ST LOC.
- Written 14 state annotations and 4 state filters.
- Median patch size is more than 10x higher than Ksplice's.
- Instrumentation cost isolated in allocator operations (1.06-2.30x).
- Instrumentation yields a modest memory overhead ($\sim 0.35x$).



Update Time



Summary

- PROTEOS: a new research OS designed with live update in mind.
- Supports several classes of updates with minimal manual effort.
- Full control over the live update transaction.
- Simple and stable live update process.
- Automated and extensible state transfer and state checking.
- State transfer error detection and recovery using hot rollback.



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